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## APPROXIMATION OF CAPACITIES WITH ADDITIVE MEASURES

For a space of non-additive regular measures on a metric compactum with the Prokhorov-style metric, it is shown that the problem of approximation of arbitrary measure with an additive measure on a fixed finite subspace reduces to linear optimization problem with parameters dependent on the values of the measure on a finite number of sets.

An algorithm for such an approximation, which is more efficient than the straightforward usage of simplex method, is presented.

*Key words and phrases:* Prokhorov metric, non-additive measure, approximation, compact metric space.

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### INTRODUCTION

Capacities were introduced by Choquet [1] and found numerous applications in different branches of mathematics. Spaces of upper semicontinuous capacities on compacta were systematically studied in [5]. In particular, in the latter paper functoriality of the construction of a space of capacities was proved and Prokhorov-style and Kantorovich-Rubinstein-style metrics on the set of capacities on a metric compactum were introduced. Needs of practice require that a capacity can be approximated with capacities of simpler structure or with some convenient properties.

We follow the terminology and notation of [5] and denote by  $\exp X$  the set of all non-empty closed subsets of a compactum  $X$ . We call a function  $c : \exp X \cup \{\emptyset\} \rightarrow I$  a *capacity* on a compactum  $X$  if the three following properties hold for all subsets  $F, G \subset X$ :

1.  $c(\emptyset) = 0$ ;
2. if  $F \subset G$ , then  $c(F) \leq c(G)$  (monotonicity);
3. if  $c(F) < a$ , then there is an open subset  $U \supset F$  such that for all  $G \subset U$  the inequality  $c(G) < a$  is valid (upper semicontinuity).

If, additionally,  $c(X) = 1$  (or  $c(X) \leq 1$ ) holds, then the capacity is called *normalized* (resp. *subnormalized*). We denote by  $\bar{M}X$ ,  $MX$ , and  $\underline{M}X$  the sets of all capacities on  $X$ , of all normalized, and of all subnormalized capacities on  $X$  respectively.

It was shown in [5] that  $MX$  carries a compact Hausdorff topology with the subbase of all sets of the form

$$O_-(F, a) = \{c \in MX \mid c(F) < a\}, \text{ where } F \subset X, a \in I,$$

and

$$\begin{aligned} O_+(U, a) &= \{c \in MX \mid c(U) > a\} \\ &= \{c \in MX \mid \text{there is a compactum } F \subset U, c(F) > a\}, \text{ where } U \underset{\text{op}}{\subset} X, a \in I. \end{aligned}$$

The same formulae determine a subbase of a compact Hausdorff topology on  $\underline{MX}$  so that  $MX \subset \underline{MX}$  is a subspace.

Previously we have considered the following subclasses of  $MX$ :

1)  $M_{\cap}X$  is the set of the so-called  $\cap$ -capacities (or necessity measures) with the property:  $c(A \cap B) = \min\{c(A), c(B)\}$  for all  $A, B \underset{\text{cl}}{\subset} X$ .

2)  $M_{\cup}X$  is the set of the so-called  $\cup$ -capacities (or possibility measures) with the property:  $c(A \cup B) = \max\{c(A), c(B)\}$  for all  $A, B \underset{\text{cl}}{\subset} X$ .

3) Class  $MX_0$  of capacities defined on a closed subspace  $X_0 \subset X$ . We regard each capacity  $c_0$  on  $X_0$  as a capacity on  $X$  extended with the formula  $c(F) = c_0(F \cap X_0)$ ,  $F \underset{\text{cl}}{\subset} X$ .

4) Class  $M_{Lip}X$  of capacities that are non-expanding w.r.t. the Hausdorff metric on  $\exp X$ .

Analogous subclasses are defined in  $\underline{MX}$  and  $\overline{MX}$ , with the obvious denotations.

It was proved in [2, 3] that the subsets  $M_{\cap}X$ ,  $M_{\cup}X$ ,  $M_{Lip}X$ , and  $MX_0$  are closed in  $MX$ , hence for a compactum  $X$  they are compacta as well, similarly for the respective subsets in  $\underline{MX}$  and  $\overline{MX}$ .

We consider the metric on the set  $\overline{MX}$  of capacities on a metric compactum  $(X, d)$  :

$$\hat{d}(c, c') = \inf\{\varepsilon > 0 \mid c(\overline{O}_{\varepsilon}(F)) + \varepsilon \geq c'(F), c'(\overline{O}_{\varepsilon}(F)) + \varepsilon \geq c(F), \forall F \underset{\text{cl}}{\subset} X\},$$

here  $\overline{O}_{\varepsilon}(F)$  is the closed  $\varepsilon$ -neighborhood of a subset  $F \subset X$ . The restrictions of this metric on  $\underline{MX}$  and  $MX$  are admissible [5].

For an arbitrary capacity  $c$  on a metric compactum  $X$ , explicit constructions for the closest to  $c$  point in the four above subclasses were presented in [3, 4].

Now we consider probably the most important class of *additive* regular measures.

Our goal is to approximate a capacity  $c$  on a metric compactum  $X$  with an additive measure on a *finite subspace* of  $X$ . Such measures are dense in the space  $\overline{PX}$  of all finite additive regular measures and have nice representation as linear combinations of Dirac measures.

## 1 ALGORITHM FOR APPROXIMATION OF A CAPACITY WITH AN ADDITIVE MEASURE ON A FINITE SUBSPACE

Consider a capacity  $c$  on a metric compactum  $(X, d)$  and a finite subspace  $X_0 = \{x_1, x_2, \dots, x_n\} \subset X$ . We are going to find the distance between  $c \in \overline{MX}$  and the subspace  $\overline{PX}_0 \subset \overline{MX}$ , in particular to find an additive measure  $m$  on  $X_0$  that is (almost) the closest to  $c$  with respect to the distance  $\hat{d}$ .

The inequality  $\hat{d}(c, m) \leq \varepsilon$  means that there is  $0 \leq z \leq \varepsilon$  satisfying

$$\begin{cases} m(A) \leq c(\overline{O}_{\varepsilon} A) + z, \\ c(A) \leq m(\overline{O}_{\varepsilon} A) + z \end{cases}$$

for all  $A \subset X$ . Obviously it is sufficient to verify the first inequality  $m(A) \leq c_{\varepsilon}^+(A) + z$ , where we denote  $c_{\varepsilon}^+ = c(\overline{O}_{\varepsilon}(A))$ , only for all  $A \subset X_0$ . Similarly, for the second condition we verify

$c(B) \leq m(A) + z$  for all  $B \subset X$  and  $A \subset X_0$  such that  $(\bar{O}_\varepsilon B) \cap X_0 \subset A$ . This is equivalent to  $m(A) \geq c_\varepsilon^-(A) - z$  for all  $A \subset X_0$ , where

$$c_\varepsilon^-(A) = c(X \setminus \bar{O}_\varepsilon(X_0 \setminus A)) = \sup_{\text{cl}} \{c(B) \mid B \subset X, B \cap \bar{O}_\varepsilon(X_0 \setminus A) = \emptyset\}.$$

Obviously  $c_\varepsilon^-(A) \leq c_\varepsilon^+(A)$  for all  $A \subset X_0$ .

All additive measures on  $X_0$  are of the form  $m = y_1\delta_{x_1} + y_2\delta_{x_2} + \dots + y_n\delta_{x_n}$ . Thus, to find the least  $z$  that satisfies the above conditions for some  $m$ , we have to solve the linear programming problem w.r.t. the variables  $y_1, y_2, \dots, y_n, z \geq 0$ :

$$\begin{cases} y_1, y_2, \dots, y_n, z \geq 0, \\ \sum_{x_i \in A} y_i \leq c_\varepsilon^+(A) + z & \text{for all } A \subset X_0, \\ \sum_{x_i \in A} y_i \geq c_\varepsilon^-(A) - z & \text{for all } A \subset X_0, \\ z \rightarrow \min, \end{cases}$$

which we rewrite as follows:

$$\begin{cases} y_1, y_2, \dots, y_n, z \geq 0, \\ -\sum_{x_i \in A} y_i + z \geq -c_\varepsilon^+(A) & \text{for all } A \subset X_0, \\ \sum_{x_i \in A} y_i + z \geq c_\varepsilon^-(A) & \text{for all } A \subset X_0, \\ z \rightarrow \min. \end{cases}$$

We embed the set  $\text{Exp } X_0$  into  $\mathbb{R}^n$  by identifying each subset  $A \subset X_0$  with the vector containing 1 at all  $i$ -th positions such that  $x_i \in A$  and 0 at all other positions. E.g.,  $\emptyset$  is represented by  $(0, \dots, 0)$ , and  $X_0$  by  $(1, \dots, 1)$ . By  $-\text{Exp } X_0$  we denote the set of the opposites to elements of  $\text{Exp } X_0 \subset \mathbb{R}^n$ . Define a function  $c_\varepsilon : \text{Exp } X_0 \cup (-\text{Exp } X_0) \rightarrow \mathbb{R}$  by the formula

$$c_\varepsilon(A) = \begin{cases} c_\varepsilon^-(A), & A \in \text{Exp } X_0, \\ -c_\varepsilon^+(-A), & A \in (-\text{Exp } X_0). \end{cases}$$

The common element  $\emptyset = (0, \dots, 0) \in \text{Exp } X_0 \cap (-\text{Exp } X_0)$  leads to no contradiction because  $c_\varepsilon^-(\emptyset) = c_\varepsilon^+(\emptyset) = 0$ .

We also denote by  $(A|1)$  the vector obtained by appending a trailing 1 to the sequence  $A = (a_1, a_2, \dots, a_n) \in \text{Exp } X_0 \cup (-\text{Exp } X_0)$ . Then the linear optimization problem can be written as

$$\begin{cases} y_1, y_2, \dots, y_n, z \geq 0, \\ (A|1) \cdot (y_1, y_2, \dots, y_n, z) \geq c_\varepsilon(A) & \text{for all } A \in \text{Exp } X_0 \cup (-\text{Exp } X_0), \\ z \rightarrow \min. \end{cases}$$

It has a straightforward geometric interpretation: of all functionals of the form

$$\gamma(t_1, t_2, \dots, t_n) = y_1 t_1 + y_2 t_2 + \dots + y_n t_n + z$$

such that  $\gamma(A) \geq c_\varepsilon(A)$  for all  $A \in \text{Exp } X_0 \cup (-\text{Exp } X_0)$ , choose one with the minimal  $z$ , i.e., with the least value  $\gamma(\vec{0})$ . Now it is clear that, due to monotonicity of the function  $c_\varepsilon$ , the restrictions  $y_1, y_2, \dots, y_n \geq 0$  can be dropped. Observe also that the restriction  $z \geq 0$  is equivalent to

$$(\emptyset|1) \cdot (y_1, y_2, \dots, y_n, z) \geq c_\varepsilon(\emptyset),$$

hence can be dropped as well.

Geometric arguments also show that the problem is solved if affinely independent

$$A_1, A_2, \dots, A_{n+1} \in \text{Exp } X_0 \cup (-\text{Exp } X_0)$$

are found such that  $\vec{0}$  is in their convex hull (in the sequel we call such  $A_1, A_2, \dots, A_{n+1}$  *basic subsets*), and the solutions  $y_1, y_2, \dots, y_n, z$  of the system

$$\begin{cases} (A_1|1) \cdot (y_1, y_2, \dots, y_n, z) & = c_\varepsilon(A_1), \\ (A_2|1) \cdot (y_1, y_2, \dots, y_n, z) & = c_\varepsilon(A_2), \\ \dots & \\ (A_{n+1}|1) \cdot (y_1, y_2, \dots, y_n, z) & = c_\varepsilon(A_{n+1}) \end{cases}$$

satisfy

$$(A|1) \cdot (y_1, y_2, \dots, y_n, z) \geq c_\varepsilon(A)$$

for all  $A \in \text{Exp } X_0 \cup (-\text{Exp } X_0)$ .

Therefore we propose the following algorithm, which essentially is equivalent to the simplex algorithm, but is better suited for our needs. Choose initial basic subsets, e.g.,  $A_1 = \{x_1\}$ ,  $A_2 = \{x_2\}$ , ...,  $A_n = \{x_n\}$ ,  $A_{n+1} = -\{x_n\}$ , then calculate  $y_1, y_2, \dots, y_n, z$  as

$$(y_1, y_2, \dots, y_n, z)^T = (M(A_1, A_2, \dots, A_n))^{-1}(c(A_1), c(A_2), \dots, c(A_{n+1}))^T,$$

where  $(-)^T$  means transposition, and

$$M(A_1, A_2, \dots, A_n) = \begin{bmatrix} A_1 & | & 1 \\ A_2 & | & 1 \\ \dots & & \dots \\ A_{n+1} & | & 1 \end{bmatrix},$$

i.e., it is the matrix with the rows  $(A_1|1), (A_2|1), \dots, (A_{n+1}|1)$ .

We will permanently need the inverse matrix

$$(M(A_1, A_2, \dots, A_n))^{-1} = \begin{bmatrix} \lambda_{11} & \lambda_{12} & \dots & \lambda_{1,n+1} \\ \lambda_{21} & \lambda_{22} & \dots & \lambda_{2,n+1} \\ \dots & \dots & \ddots & \dots \\ \lambda_{n1} & \lambda_{n2} & \dots & \lambda_{n,n+1} \\ \mu_1 & \mu_2 & \dots & \mu_{n+1} \end{bmatrix}.$$

For any  $A \in \text{Exp } X_0 \cup (-\text{Exp } X_0)$  the column  $(M(A_1, A_2, \dots, A_n))^{-1}(A|1)^T$  consists of the coefficients  $\alpha_1, \alpha_2, \dots, \alpha_{n+1}$  such that  $\alpha_1 + \alpha_2 + \dots + \alpha_{n+1} = 1$  and  $\alpha_1 A_1 + \alpha_2 A_2 + \dots + \alpha_{n+1} A_{n+1} = A$  (in the above sense). In particular,  $\mu_1 A_1 + \mu_2 A_2 + \dots + \mu_{n+1} A_{n+1} = \emptyset$ , and  $\lambda_{i1} A_1 + \lambda_{i2} A_2 + \dots + \lambda_{i,n+1} A_{n+1} = \{x_i\}$  for all  $1 \leq i \leq n$ .

Now, having  $y_1, y_2, \dots, y_n, z$  calculated, compare the differences

$$c_\varepsilon(A) - (A|1)(y_1, y_2, \dots, y_n, z)$$

for all  $A \in \text{Exp } X_0 \cup (-\text{Exp } X_0)$ . If the basic subsets  $A_1, A_2, \dots, A_{n+1}$  provide a solution, then all the differences are not greater than 0. Otherwise find the greatest difference  $\Delta =$

$c_\varepsilon(A') - (A'|1)(y_1, y_2, \dots, y_n, z)$ , which is positive, and replace with  $A'$  a subset  $A_i$  such that  $\vec{0}$  is in the convex hull of  $A_1, A_2, \dots, A_{i-1}, A', A_{i+1}, \dots, A_{n+1}$ .

Let  $(\alpha_1, \alpha_2, \dots, \alpha_{n+1})^T = (M(A_1, A_2, \dots, A_n))^{-1}(A'|1)^T$ , hence  $A' = \alpha_1 A_1 + \alpha_2 A_2 + \dots + \alpha_{n+1} A_{n+1}$ , then

$$A_i = \frac{1}{\alpha_i} A' - \frac{\alpha_1}{\alpha_i} A_1 - \dots - \frac{\alpha_{i-1}}{\alpha_i} A_{i-1} - \frac{\alpha_{i+1}}{\alpha_i} A_{i+1} - \frac{\alpha_{n+1}}{\alpha_i} A_{n+1}.$$

Therefore

$$\begin{aligned} \emptyset &= (\mu_1 - \mu_i \frac{\alpha_1}{\alpha_i}) A_1 + \dots + (\mu_{i-1} - \mu_i \frac{\alpha_{i-1}}{\alpha_i}) A_{i-1} + (\mu_{i+1} - \mu_i \frac{\alpha_{i+1}}{\alpha_i}) A_{i+1} \\ &+ \dots + (\mu_{n+1} - \mu_i \frac{\alpha_{n+1}}{\alpha_i}) A_{n+1} + \frac{\mu_i}{\alpha_i} A'. \end{aligned}$$

The coefficients in the new decomposition of  $\emptyset$  should be nonnegative, hence  $\alpha_i > 0$  is required, as well as either  $\alpha_j \leq 0$  or  $\mu_j - \mu_i \frac{\alpha_j}{\alpha_i} \geq 0$  for all  $j \neq i$ . If  $\alpha_j > 0$ , then the latter inequality is equivalent to  $\frac{\mu_j}{\alpha_j} \geq \frac{\mu_i}{\alpha_i}$ . Hence  $\frac{\mu_i}{\alpha_i}$  should be the least of  $\frac{\mu_j}{\alpha_j}$  for  $1 \leq j \leq n+1$  such that  $\alpha_j > 0$ .

Now we replace  $A_i$  with  $A'_i = A'$ , and the inverse matrix

$$(M(A_1, A_2, \dots, A_{i-1}, A'_i, A_{i+1}, \dots, A_n))^{-1} = \begin{bmatrix} \lambda'_{11} & \lambda'_{12} & \dots & \lambda'_{1,n+1} \\ \lambda'_{21} & \lambda'_{22} & \dots & \lambda'_{2,n+1} \\ \dots & \dots & \ddots & \dots \\ \lambda'_{n1} & \lambda'_{n2} & \dots & \lambda'_{n,n+1} \\ \mu'_1 & \mu'_2 & \dots & \mu'_{n+1} \end{bmatrix}$$

is adjusted accordingly:

$$\begin{aligned} \mu'_i &= \frac{\mu_i}{\alpha_i}, & \mu'_j &= \mu_j - \alpha_j \frac{\mu_i}{\alpha_i}, & 1 \leq j \leq n+1, j \neq i, \\ \lambda'_{ki} &= \frac{\lambda_{ki}}{\alpha_i}, & \lambda'_{kj} &= \lambda_{kj} - \alpha_j \frac{\lambda_{ki}}{\alpha_i}, & 1 \leq k, j \leq n+1, j \neq i. \end{aligned}$$

Now look how  $y_1, y_2, \dots, y_n, z$  have changed. Taking into account

$$\begin{aligned} z &= \mu_1 c_\varepsilon(A_1) + \dots + \mu_{i-1} c_\varepsilon(A_{i-1}) + \mu_i c_\varepsilon(A_i) \\ &+ \mu_{i+1} c_\varepsilon(A_{i+1}) + \dots + \mu_{n+1} c_\varepsilon(A_{n+1}), \\ z' &= (\mu_1 - \alpha_1 \frac{\mu_i}{\alpha_i}) c_\varepsilon(A_1) + \dots + (\mu_{i-1} - \alpha_{i-1} \frac{\mu_i}{\alpha_i}) c_\varepsilon(A_{i-1}) + \frac{\mu_i}{\alpha_i} c_\varepsilon(A'_i) \\ &+ (\mu_{i+1} - \alpha_{i+1} \frac{\mu_i}{\alpha_i}) c_\varepsilon(A_{i+1}) + \dots + (\mu_{n+1} - \alpha_{n+1} \frac{\mu_i}{\alpha_i}) c_\varepsilon(A_{n+1}), \end{aligned}$$

obtain

$$z' - z = \frac{\mu_i}{\alpha_i} (c_\varepsilon(A'_i) - (\alpha_1 c_\varepsilon(A_1) + \dots + \alpha_{n+1} c_\varepsilon(A_{n+1}))) = \frac{\mu_i}{\alpha_i} \cdot \Delta.$$

Similarly

$$y'_k - y_k = \frac{\lambda_{ki}}{\alpha_i} (c_\varepsilon(A'_i) - (\alpha_1 c_\varepsilon(A_1) + \dots + \alpha_{n+1} c_\varepsilon(A_{n+1}))) = \frac{\lambda_{ki}}{\alpha_i} \cdot \Delta.$$

This simplifies calculation of  $z'$  and all  $y'_k$ . We iterate the above step until  $\Delta = 0$ . The final value of  $z$ , which we denote  $z(\varepsilon)$ , is the least  $z$  such that

$$\begin{cases} m(A) \leq c(\bar{O}_\varepsilon A) + z, \\ c(A) \leq m(\bar{O}_\varepsilon A) + z \end{cases}$$

for some  $m \in \bar{P}X_0$  and all  $A \subset_{\text{cl}} X$ .

Observe that  $z(\varepsilon)$  is non-increasing with respect to  $\varepsilon$ , hence the distance between  $c$  and  $\bar{P}X_0$  is the least  $\varepsilon$  such that  $z(\varepsilon) \leq \varepsilon$ . This distance is not greater than  $z(0)$ , therefore it is easy to bisect the segment  $[0, z(0)]$  to find the distance and an approximating additive measure with arbitrary precision.

## 2 CONCLUDING REMARKS

The proposed algorithm was implemented as a C program and tested on data sets with cardinality of  $X_0$  up to 10.

However, each iteration of the presented algorithm requires previously calculated values of a capacity for all  $2^{\text{cardinality of the space}}$  subsets, which is not appropriate even for  $\geq 40$  points. Hence, to handle subspaces of greater cardinality, we need to cut memory and time requirements using the metric structure and the only reliable property of a capacity, i.e., its monotonicity. This requires deeper investigation combining both topological properties of non-additive measures, e.g., their dimensional characteristics, and computational aspects.

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Для простору неадитивних регулярних мір на метричному компактi з відстанню в стилі Прохорова показано, що задача наближення довільної міри адитивною мірою на фіксованому скінченному підпросторі зводиться до задачі лінійної оптимізації з параметрами, залежними від значень вихідної міри на скінченному числі множин.

Запропоновано алгоритм такого наближення, ефективніший порівняно з прямолінійним застосуванням симплекс-методу.

*Ключові слова і фрази:* метрика Прохорова, неадитивна міра, апроксимація, компактний метричний простір.